# Bucknell SUniversity 

## CSCI 311 - Data Structures

The Knapsack Problem

## The Knapsack Problem

Knapsack of capacity M


N Types of Indivisible Items (unlimited number of each type)


Problem: What is the selection of items that fits in the knapsack maximizing the total value of its contents?

## The Knapsack Problem

|  |  | Type A | Type B | Type C |
| :--- | :--- | :---: | :---: | :---: |
| $\mathrm{N}=3$ <br> $\mathrm{M}=8$ | Value | 100 | 76 | 54 |
|  | Weight | 5 | 4 | 3 |



## Note:

- For each node in this tree, we have a set of possible decisions.
- Each decision has a cost (its weight) and leads to an associated yield.
- The goal is to find a sequence of decisions that leads to an optimal solution.
- The number of possible solutions is exponential with M. We'd have to find them all and then choose the very best.


## The Knapsack Problem

The recursive nature of the problem jumps out at us when we observe the decision tree.

The problem has optimal substructure and overlapping sub-problems, so it is solvable with dynamic programming.

What we have to figure out is how to map the problem onto some kind of data structure to store solutions to each subproblem as the tree is traversed.

## The Knapsack Problem

| $\mathrm{N}=3$ <br> $\mathrm{M}=27$ |
| :--- | :--- | :---: | :---: | :---: |
|  Type A Type B Type C <br> Value 100 76 54 <br>  Weight 5 4 |



Question: What kind of data structure is needed to apply DP to this problem?

## The Knapsack Problem (recursive solution)

## knap( M)

$\max =0$;
for $\mathbf{i}=1$ to $N / /$ Loop through itemtypes
// Sol ve probl em assumi ng we i nclude // an item of type i do spaceLeft $=M$ - size[i] if spaceLeft $>=0 / /$ if type ifits then // Compute candi date sol'n t $\mathrm{t}=$ knap(spaceLeft) +val [i] if $t>\max$ then max $=t$
ret urn max;

## The Knapsack Problem (DP solution)

```
Knap(M)
    if maxKnown[M]!= unknown
        then return maxKnown[M];
    // Othervise, result not yet known:
    max = O
    for i = l to N // Try each itemtype
        do spaceLeft = M - size[i]
            if spaceLeft >=0 // lf itemtype i fits
            then // Compute candi date sol ution t
            t = knap(spaceLeft) + val [i]
            if t > max
                        then max = t;
                                maxi = i;
maxKnown[M] = max // memoizee result
return max
```

