

Binary Search Tree

Revised based on textbook
author's notes.

Search Trees

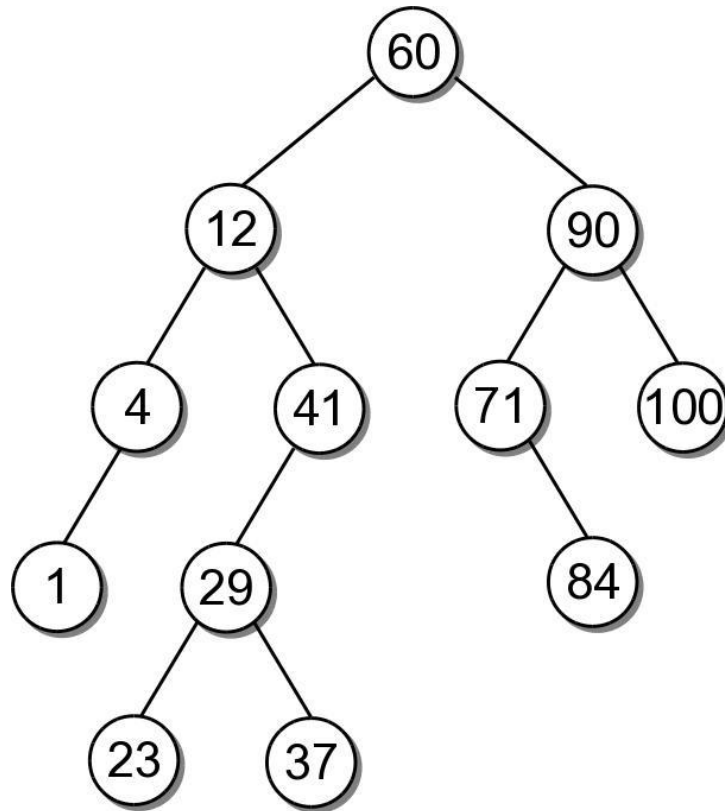
- The tree structure can be used for searching.
 - Each node contains a search key as part of its data or **payload**.
 - Nodes are organized based on the relationship between the keys.
- Search trees can be used to implement various types of containers.
 - Most common use is with the Map ADT.

Binary Search Tree (BST)

- A binary tree in which each node contains a search key and the tree is structured such that for each interior node V :
 - All keys less than the key in node V are stored in the left subtree of V .
 - All keys greater than the key in node V are stored in the right subtree of V .

BST Example

- Consider the example tree



BST – ADT

bst.py

```
# We use an unique name to distinguish this version  
# from others in the chapter.
```

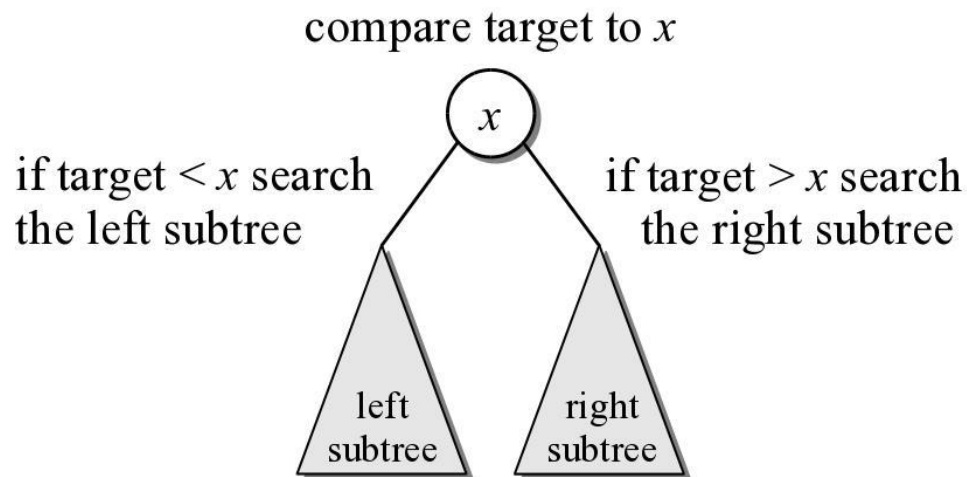
```
class BST :  
    def __init__( self ):  
        self._root = None  
        self._size = 0  
  
    def __len__( self ):  
        return self._size  
  
    def __iter__( self ):  
        return _BSTreeIterator( self._root )  
# ...
```

```
# Storage class for the binary search tree nodes.
```

```
class _BSTNode :  
    def __init__( self, key, data ):  
        self.key = key  
        self.data = data  
        self.left = None  
        self.right = None
```

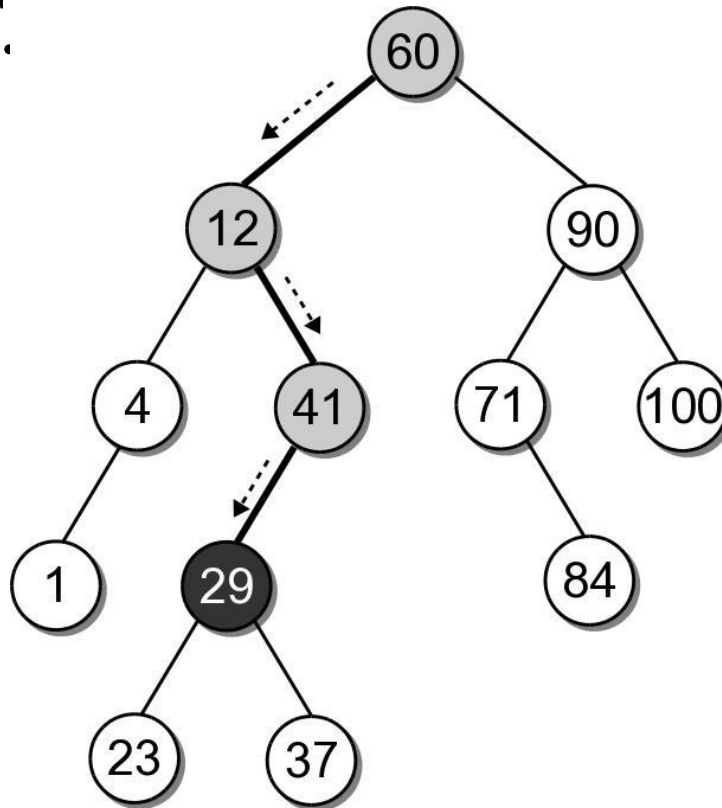
BST – Searching

- A search begins at the root node.
 - The target is compared to the key at each node.
 - The path depends on the relationship between the target and the key in the node.



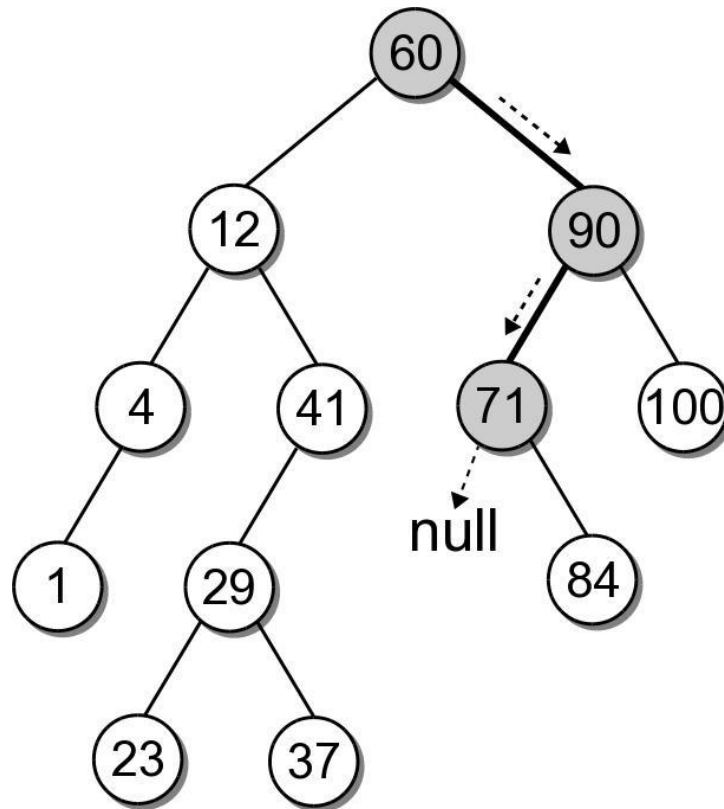
BST – Search Example

- Suppose we want to search for 29 in our BST.



BST – Search Example

- What if the key is not in the tree?
Search for key 68 in our BST.



BST – Search Implementation

bst.py

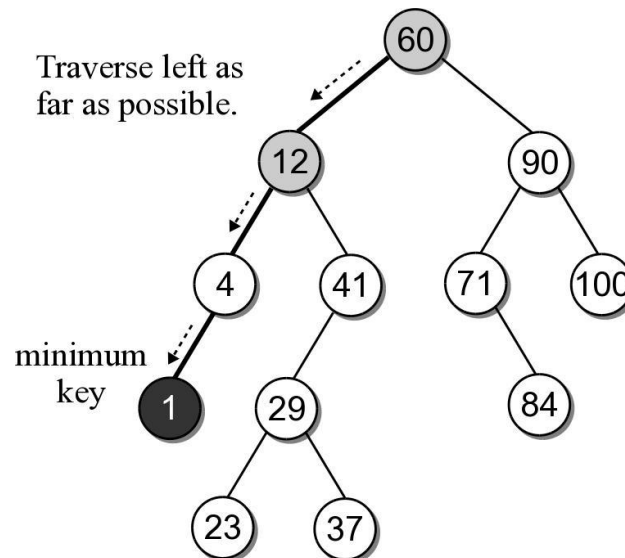
```
class BST :
# ...
    def __contains__( self, key ) :
        return self._bstSearch( self._root, key ) is not None

    def valueOf( self, key ) :
        node = self._bstSearch( self._root, key )
        assert node is not None, "Invalid map key."
        return node.value

    def _bstSearch( self, subtree, target ) :
        if subtree is None :
            return None
        elif target < subtree.key :
            return self._bstSearch( subtree.left )
        elif target > subtree.key :
            return self._bstSearch( subtree.right )
        else :
            return subtree
```

BST – Min or Max Key

- Finding the minimum or maximum key within a BST is similar to the general search.
 - Where might the smallest key be located?
 - Where might the largest key be located?



BST – Min or Max Key

- The helper method below finds the node containing the minimum key.

```
class BST :  
    # ...  
    def _bstMinumum( self, subtree ) :  
        if subtree is None :  
            return None  
        elif subtree.left is None :  
            return subtree  
        else :  
            return self._bstMinimum( subtree.left )
```

BST – Insertions

- When a BST is constructed, the keys are added one at a time. As keys are inserted
 - A new node is created for each key.
 - The node is linked into its proper position within the tree.
 - The search tree property must be maintained.

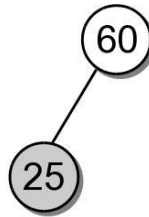
Building a BST

- Suppose we want to build a BST from the key list

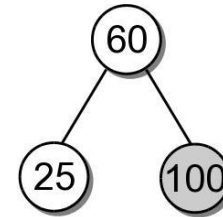
60 25 100 35 17 80



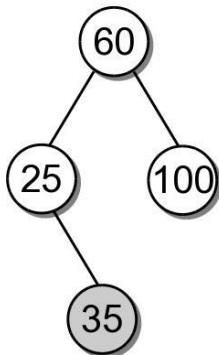
(a) Insert 60.



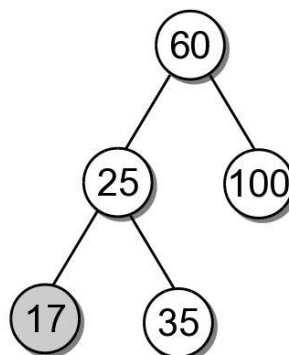
(b) Insert 25.



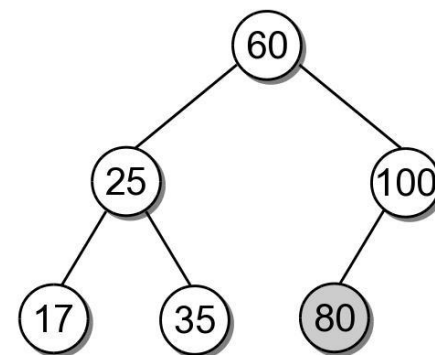
(c) Insert 100.



(d) Insert 35.



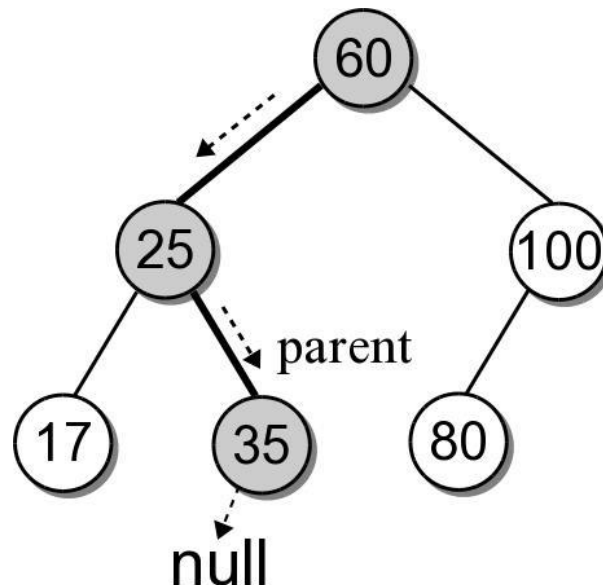
(e) Insert 17.



(f) Insert 80.

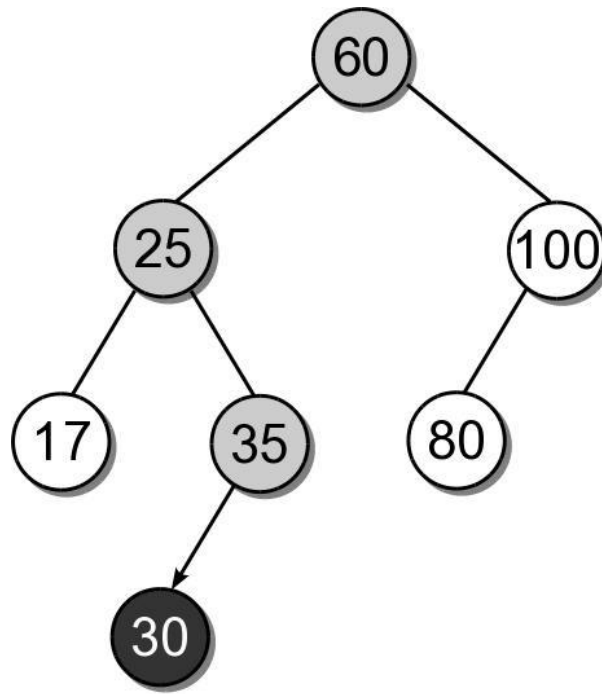
BST – Insertion

- Building a BST by hand is easy. How do we insert an entry in program code?
 - What happens if we use the search method from earlier to search for key 30?



BST – Insertion

- We can insert the new node where the search fell off the tree.



BST – Insert Implementation

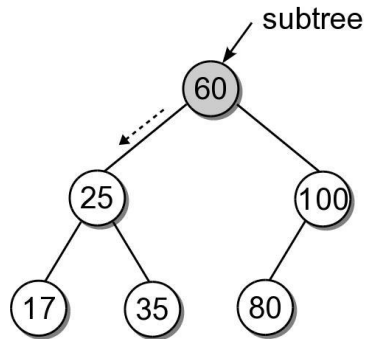
bst.py

```
class BST :
# ...
    def add( self, key, value ) :
        node = self._bstSearch( key )
        if node is not None :
            node.value = value
            return False
        else :
            self._root = self._bstInsert( self._root, key, value )
            self._size += 1
            return True

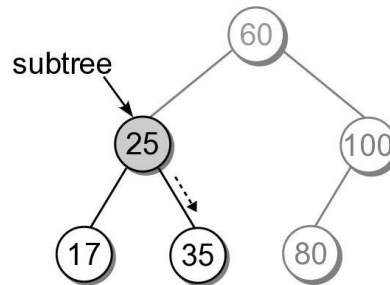
    def _bstInsert( self, subtree, key, value ) :
        if subtree is None :
            subtree = _BSTMapNode( key, value )
        elif key < subtree.key :
            subtree.left = self._bstInsert(subtree.left, key, value)
        elif key > subtree.key :
            subtree.right = self._bstInsert(subtree.right, key, value)
        return subtree
```


BST – Insert Steps

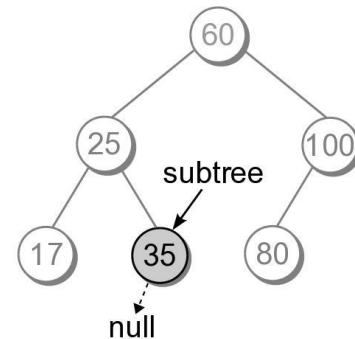
- Add 30 to our sample BST.



(a) `bstInsert(root,30)`

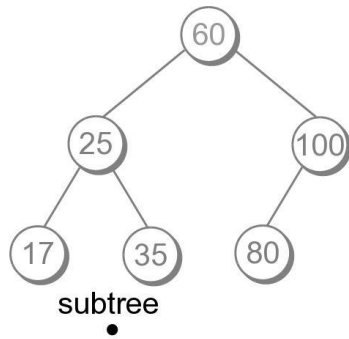


(b) `bstInsert(subtree.left,key)`

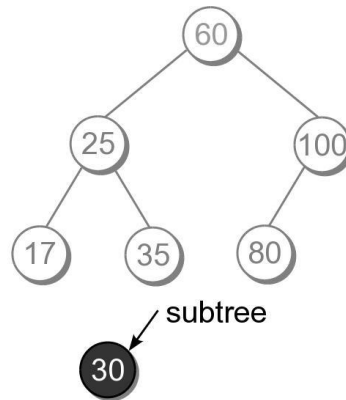


(c) `bstInsert(subtree.right,key)`

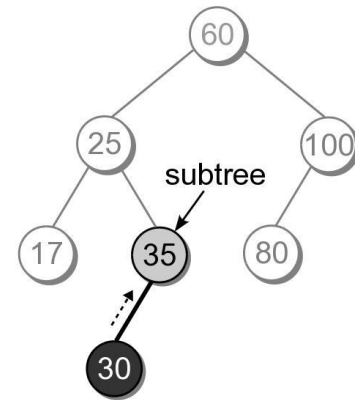
BST – Insert Steps



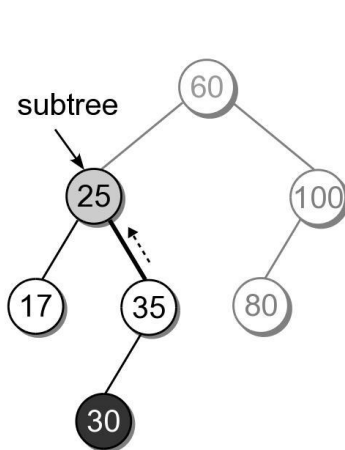
(d) `bstInsert(subtree.left, key)`



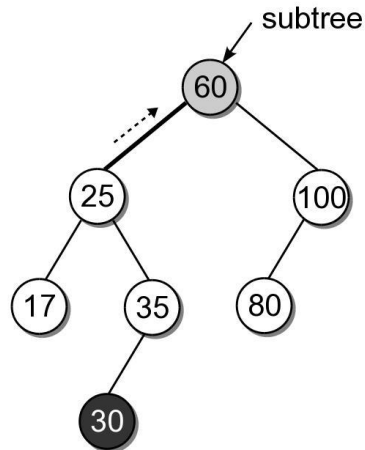
(e) `subtree = TreeNode(key)`



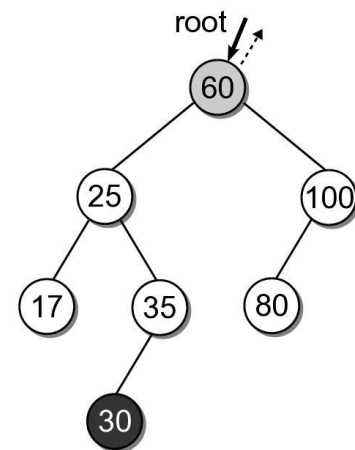
(f) `subtree.left = bstInsert(...)`



(g) `subtree.right = bstInsert(...)`



(h) `subtree.left = bstInsert(...)`



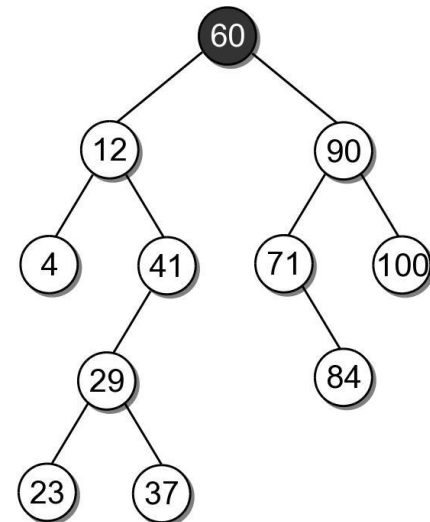
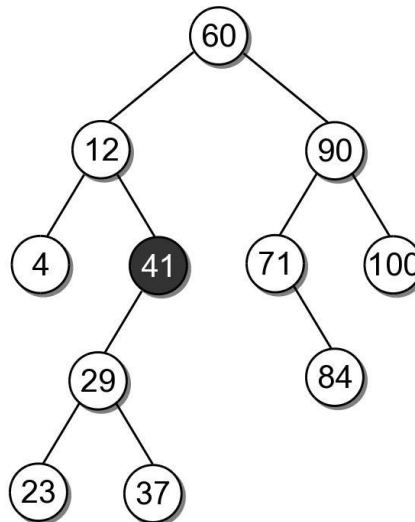
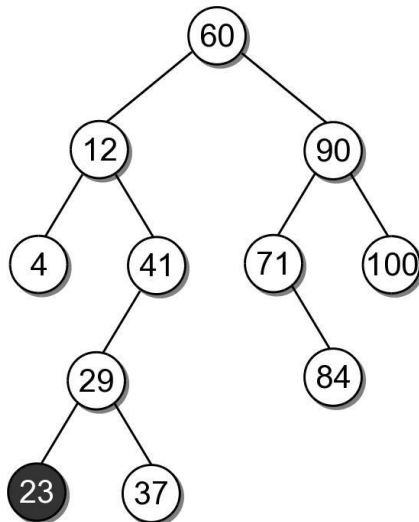
(i) `root = bstInsert(...)`

BST – Deletions

- Deleting a node from a BST is a bit more complicated.
 - Locate the node containing the node.
 - Delete the node.
- When a node is removed, the remaining nodes must preserve the search tree property.

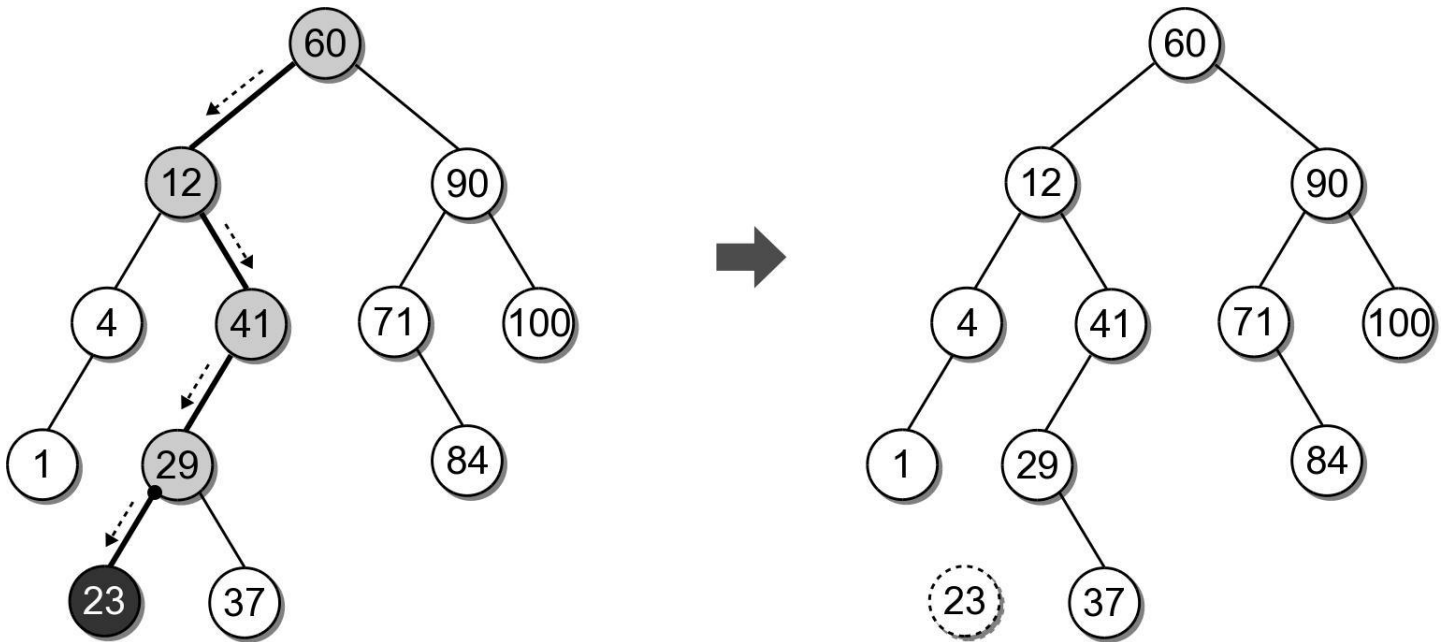
BST – Deletions

- There are three cases to consider:
 - the node is a leaf.
 - the node has a single child
 - the node has two children.



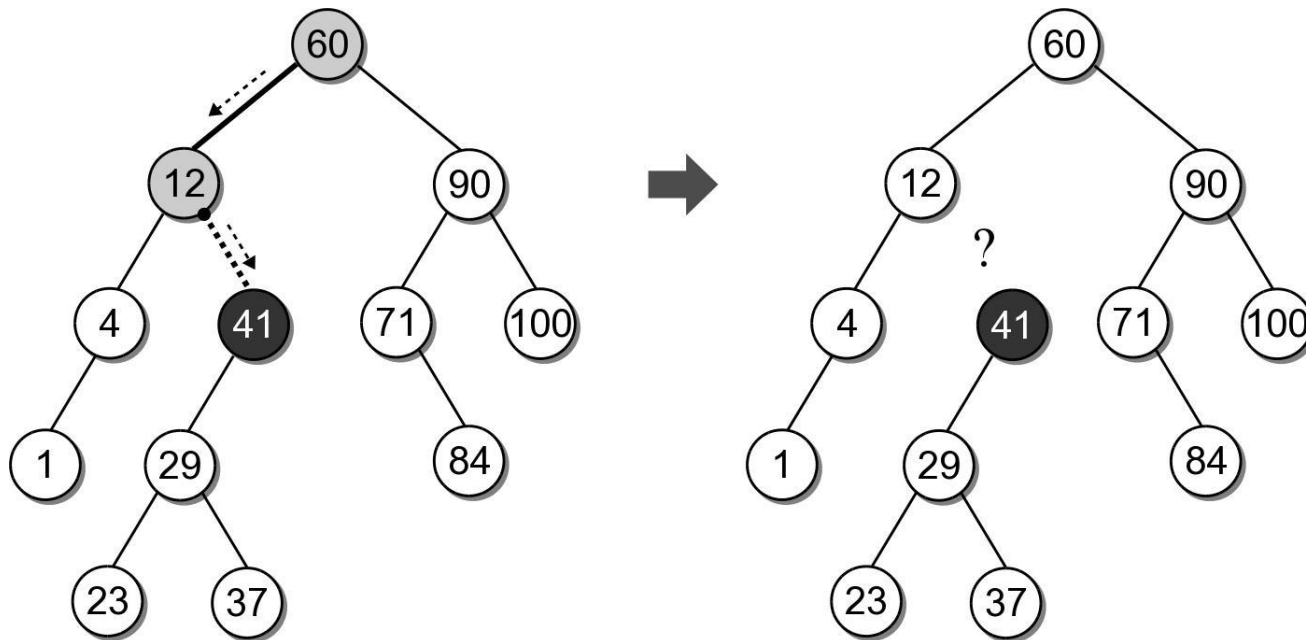
BST – Delete Leaf Node

- Removing a leaf node is the easiest case.
- Suppose we want to remove 23.



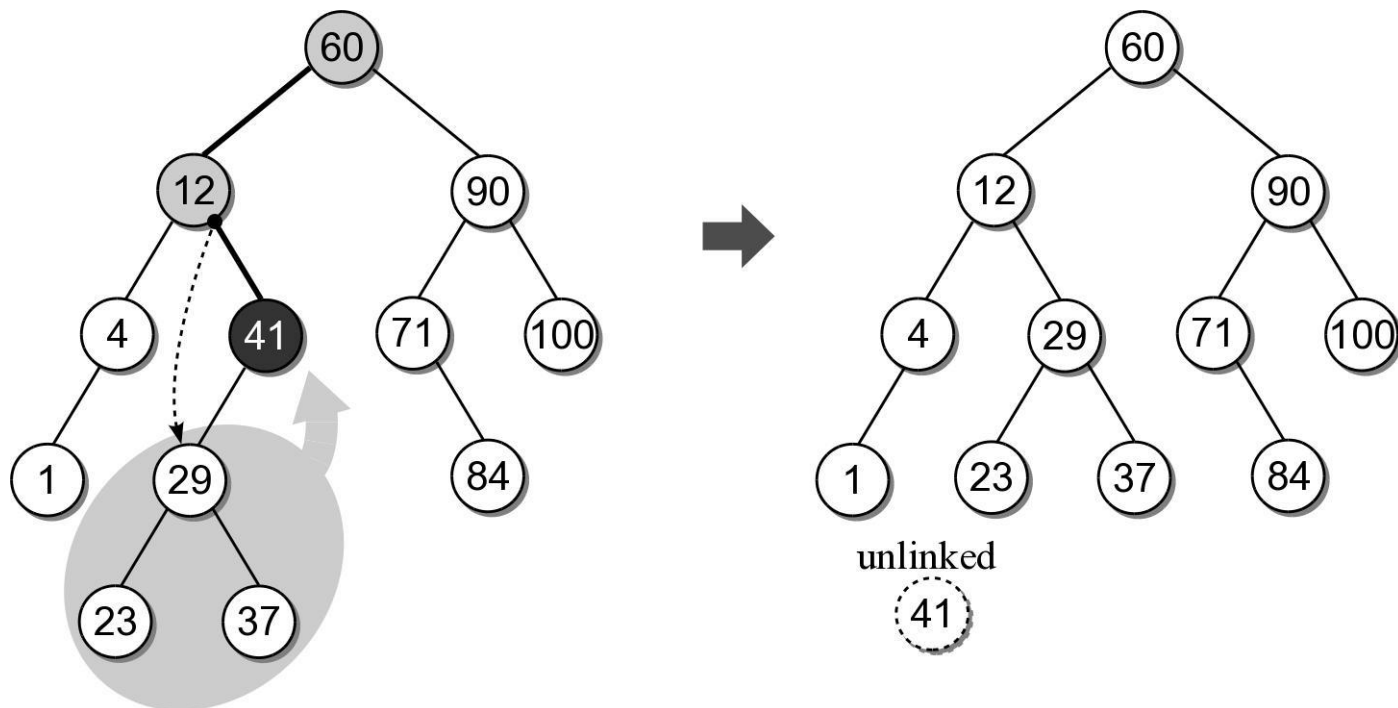
BST – Delete Interior Node

- Removing an interior node with one child.
 - Suppose we want to remove 41.
 - We can not simply unlink the node.



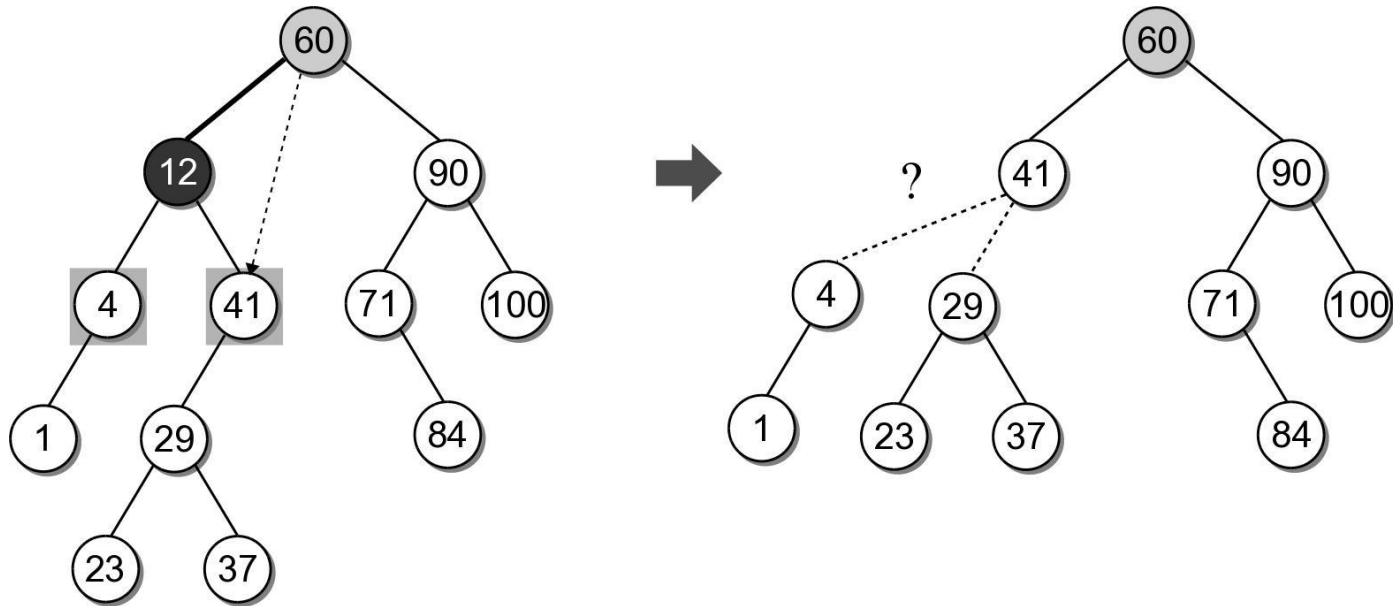
BST – Delete Interior Node

- After locating the node to be removed, it's child must be linked to it's parent.



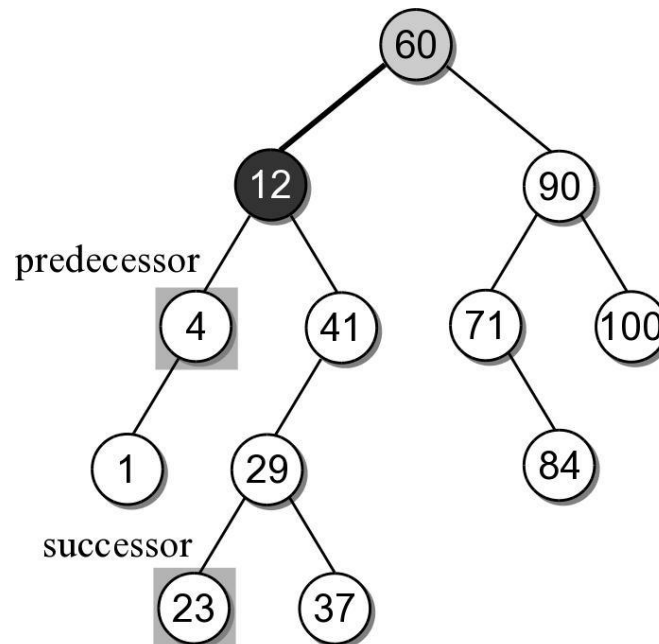
BST – Delete Interior Node

- The most difficult case is deleting a node with two children.
 - Suppose we want to delete node 12.
 - Which child should be linked to the parent?



BST – Delete Interior Node

- Based on the search tree property, each node has a logical predecessor and successor.
 - For node 12, those are 4 and 23.

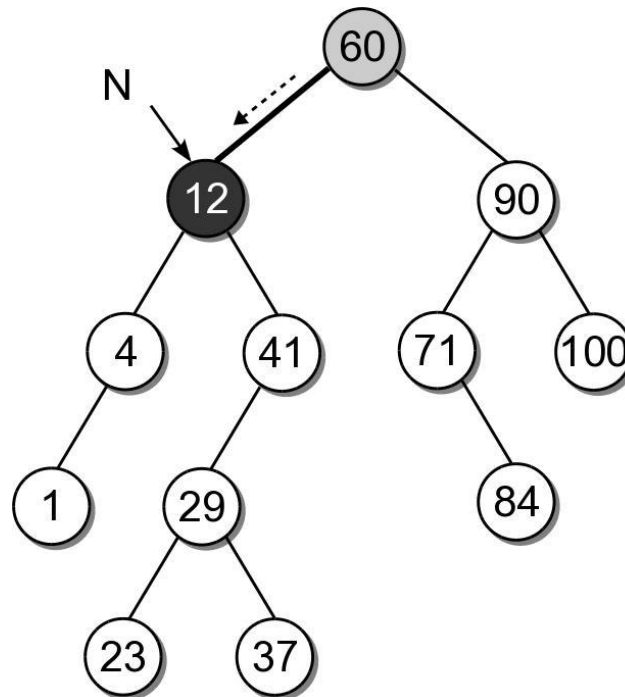


BST – Delete Interior Node

- We can replace to be deleted with either its logical successor or predecessor.
 - Both will either be a leaf or an interior node with one child.
 - We already know how to remove those nodes.

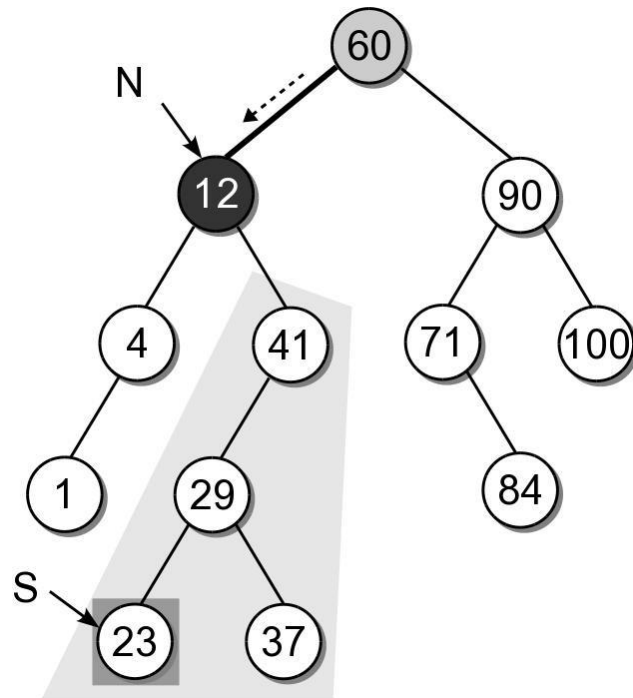
BST – Delete Interior Node

- Removing an interior node with two children requires 4 steps:
 - (1) Find the node to be deleted, N.



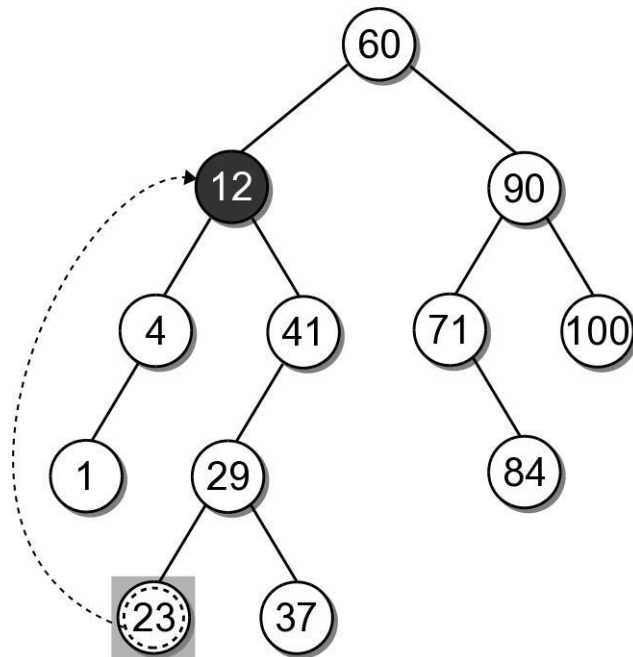
BST – Delete Interior Node

- (2) Find the successor, S, of node N.



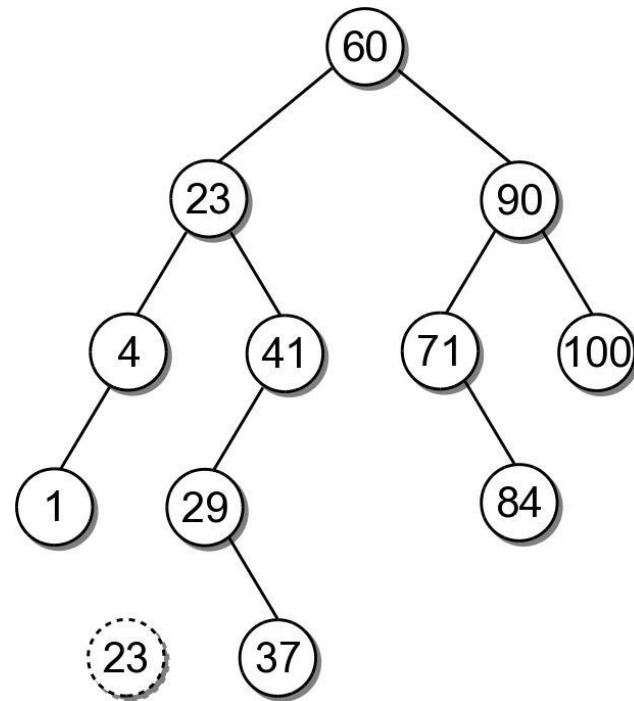
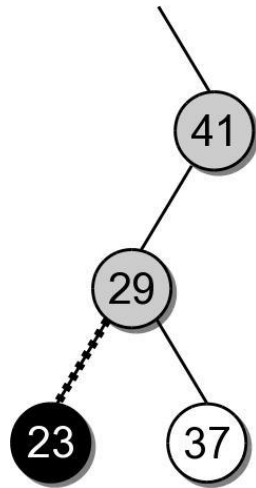
BST – Delete Interior Node

- (3) Copy the payload from node S to node N.



BST – Delete Interior Node

- (4) Remove node S from the tree.



BST – Delete Interior Node

- Removing an interior node with two children requires 4 steps:
 - Find the node to be deleted, N .
 - Find the logical successor, S , of node N .
 - Copy the payload from node S to node N .
 - Remove node S from the tree.

BST – Delete Implementation

bst.py

```
class BST :  
    # ...  
    def remove( self, key ) :  
        assert key in self, "Invalid key."  
        self._root = self._bstRemove( self._root, key )  
        self._size -= 1
```


BST – Delete Implementation

bst.py

```
class BST :  
# ...  
def _bstRemove( self, subtree, target ):  
    if subtree is None : # not found  
        return subtree  
    elif target < subtree.key : # search for the left tree  
        subtree.left = self._bstRemove( subtree.left, target )  
        return subtree  
    elif target > subtree.key : # search for the right tree  
        subtree.right = self._bstRemove( subtree.right, target )  
        return subtree  
    else : # target == subtree.key, delete the node  
        .....
```

BST – Delete Implementation

bst.py

```
class BST :
# ...
    def _bstRemove( self, subtree, target ) :
        .....
    else :
        if subtree.left is None and subtree.right is None :
            return None      # leaf node
        elif subtree.left is None or subtree.right is None :
            # has one child only
            if subtree.left is not None :
                return subtree.left
            else :
                return subtree.right
        else : # has both children
            successor = self._bstMinimum( subtree.right )
            subtree.key = successor.key
            subtree.value = successor.value
            subtree.right = self._bstRemove( subtree.right,
                                             successor.key )

        return subtree
```

BST – Efficiency

Operation	Worst Case
<code>_bstSearch(root, k)</code>	$O(n)$
<code>_bstMinimum(root)</code>	$O(n)$
<code>_bstInsert(root, k)</code>	$O(n)$
<code>_bstRemove(root, k)</code>	$O(n)$
traversal	$O(n)$

Exercises

1. The `remove()` method uses a `bst_find_min(subtree)` method to search and return the node with minimum key in a tree rooted at `subtree`. Implement this method.
2. Write a method `bst_sort_reverse()` to sort a binary search tree rooted at `subtree` in reverse order.
3. Test your methods by running the program `testbst.py` on the course website.