

# **Theory of Computation** CSCI 341, Fall 2016

## Homework 2 Due 2016-09-23

## (2pt) Problem 1

Following the spirit of the algorithm converting regular expressions to NFAs, propose an extension of this algorithm to convert regular expressions of the form  $E = E_1^+$  to NFAs. Draw the conversion **and** define it formally, by assuming  $E_1$  is recognized by an NFA  $(Q_1, \Sigma, \Delta_1, s_1, F_1)$ , define an NFA  $N = (Q, \Sigma, \Delta, s, F)$  recognizing  $E_1^+$ .

SOLUTION Let  $M_1 = (Q_1, \Sigma, \Delta_1, s_1, F_1)$  be an NFA recognizing the regexp  $E_1$ . We define the following NFA *M* recognizing  $E_1^+$  by defining it as

$$M = (Q, \Sigma, \Delta, s, F)$$

with

- $Q = Q_1$
- *s* = *s*1

• 
$$\Delta = \Delta_1 \cup \{(f, \varepsilon, s_1) \mid f \in F_1\}$$

• 
$$F = F_1$$

#### (3pt) Problem 2

Convert the following regular expressions to NFAs using the procedure we saw in class, and then convert the NFAs to DFAs using the other procedure we saw in class.

- 1.  $a^*(a(b^*))b$
- 2.  $(a \cup b)^* b(aa)b$
- 3.  $a(b^+)(a \cup b)^*(b^+)a$

#### (2.5pt) Problem 3

Exercise 4.7 from Sipser, page 211. Maximum credit for a clear and complete proof.

SOLUTION Let B be the set containing all the infinite sequences of booleans. For example, elements of B could be sequence s like:

 $s = 0000000000 \cdots$ 

$s = 11111111111\cdots$
$s = 01010101 \cdots$
$s = 0110001110\cdots$

Let us call call s[i] the *i*-th digit of *s*.

We want to prove that *B* is *uncountable*. Let's prove it by contradiction.

Let us assume that *B* is countable. It means that there exists an isomorphism  $f : \mathbb{N} \to B$ . Let us define a particular sequence *d* (the diagonal sequence) using the function *f*.

$$d[i] = \overline{f(i)[i]}$$

The function f is an bijection. Therefore, that there exists an integer k such that f(k) = d. The contradiction appears when we inspect the value of d[k]. There are two possibilities: d[k] = 0 or d[k] = 1.

- Suppose  $d[k] = 0 \Rightarrow \overline{f(k)[k]} = 0 \Rightarrow f(k)[k] = 1 \Rightarrow d[k] = 1$ . It's absurd!
- Suppose  $d[k] = 1 \Rightarrow \overline{f(k)[k]} = 1 \Rightarrow f(k)[k] = 0 \Rightarrow d[k] = 0$ . It's absurd!

This is a contradiction! In conclusion, it cannot be true that B is countable.

## (2.5pt) Problem 4

Problem 1.32 from Sipser, page 88.